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IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

Docket No. **AUS990938US1**

Assistant Commissioner for Patents
Washington, D.C. 20231

Sir:
Transmitted herewith for filing is the patent application of Inventor(s):

RICHARD LOUIS ARNDT

For: **HYPERVISOR VIRTUALIZATION OF OS CONSOLE AND OPERATOR PANEL**

Enclosed are also:

X	20	Pages of Specification including an Abstract
X	6	Pages of Claims
X	**9**	Sheet(s) of Drawings
X		A Declaration and Power of Attorney
X		Form PTO 1595 and assignment of the invention to IBM Corporation

CLAIMS AS FILED

FOR	Number Filed		Number Extra		Rate		Basic Fee ($690)
Total Claims	24	-20 =	4	X	$ 18	=	$ 72.00
Independent Claims	5	-3 =	2	X	$ 78	=	$ 156.00
Multiple Dependent Claims	0			X	$260	=	$ 0.00
					Total Filing Fee	=	$ 918.00

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X Any patent application processing fees under 37CFR § 1.17.

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Docket No. AUS990938US1

HYPERVISOR VIRTUALIZATION OF OS CONSOLE AND OPERATOR PANEL

BACKGROUND OF THE INVENTION

1. Technical Field:

The present invention relates generally to the field of computer architecture and, more specifically, to methods and systems for allowing multiple operating system images within a logically partitioned data processing system to interact with a console and operator panel.

2. Description of Related Art:

A logical partitioning option (LPAR) within a data processing system allows multiple copies of a single operating system (OS) or multiple heterogeneous operating systems to be simultaneously run on a single data processing system platform. A partition, within which an operating system image runs, is assigned a non-overlapping sub-set of the platform's resources. These platform allocable resources include one or more architecturally distinct processors with their interrupt management area, regions of system memory, and I/O adapter bus slots. The partition's resources are represented by its own open firmware device tree to the OS image.

Each distinct OS or image of an OS running within the platform are protected from each other such that software errors on one logical partition cannot affect the correct operation of any of the other partitions.

This is provided by allocating a disjoint set of platform resources to be directly managed by each OS image and by providing mechanisms for ensuring that the various images cannot control any resources that have not been allocated to it. Furthermore, software errors in the control of an OS's allocated resources are prevented from affecting the resources of any other image. Thus, each image of the OS (or each different OS) directly controls a distinct set of allocable resources within the platform.

There are certain resources within many server platforms that exist singly, yet each distinct OS within the platform must interact with these resources. For example, the RS/6000, a product of International Business Machines Corporation of Armonk, New York, includes a console and an operator panel for allowing a system administrator to detect and correct problems within the platform. However, each of these resources exists singly within the platform and it is impractical to duplicate these resources. While present architectures often do not preclude the sharing of allocable resources of this type between partitions, there is no current architectural support for such sharing. Therefore, a method, system, and computer program product for providing the sharing of allocable resources within a logically partitioned platform is desirable.

SUMMARY OF THE INVENTION

The present invention provides a logically partitioned data processing system in which shared resources are emulated to provide each partition a separate copy of the shared resource. In one embodiment, the logically partitioned data processing system includes a plurality of logical partitions, a plurality of operating systems executing within the data processing system and a plurality of assignable resources. Each of the plurality of operating systems is assigned to a separate one of the plurality of logical partitions, such that no more than one operating system is assigned to any given logical partition. Each of the plurality of assignable resources is assigned to a single one of the plurality of logical partitions. The logically partitioned data processing system also includes a hypervisor. The hypervisor emulates shared resources, such as an operator panel and a system console, and provides a virtual copy of these shared resources to each of the plurality of logical partitions.

BRIEF DESCRIPTION OF THE DRAWINGS

The novel features believed characteristic of the invention are set forth in the appended claims. The invention itself, however, as well as a preferred mode of use, further objectives and advantages thereof, will best be understood by reference to the following detailed description of an illustrative embodiment when read in conjunction with the accompanying drawings, wherein:

Figure 1 depicts a pictorial representation of a distributed data processing system in which the present invention may be implemented;

Figure 2, a block diagram of a data processing system in accordance with the present invention is illustrated;

Figure 3 depicts a block diagram of a data processing system, which may be implemented as a logically partitioned server, in accordance with the present invention;

Figure 4 depicts a block diagram illustrating a prior art logically partitioned platform in accordance with the present invention;

Figure 5 depicts a block diagram of a logically partitioned platform in which the present invention may be implemented;

Figures 6A-6B depict high-level flowcharts illustrating exemplary processes, performed for example, in hypervisor **510**, for emulating a console and operator platform in accordance with the present invention;

Figure 7 depicts a high level flowchart illustrating an exemplary process on a hardware system console for presenting the information from the various OS images to an operator in accordance with the present invention; and

Figure 8 depicts a high level flowchart illustrating an exemplary process on a hardware system console for sending messages to various ones of multiple OS images running on a logically partitioned platform in accordance with the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

With reference now to the figures, and in particular with reference to **Figure 1**, a pictorial representation of a distributed data processing system is depicted in which the present invention may be implemented.

Distributed data processing system **100** is a network of computers in which the present invention may be implemented. Distributed data processing system **100** contains network **102**, which is the medium used to provide communications links between various devices and computers connected within distributed data processing system **100**. Network **102** may include permanent connections, such as wire or fiber optic cables, or temporary connections made through telephone connections.

In the depicted example, server **104** is connected to hardware system console **150**. Server **104** is also connected to network **102**, along with storage unit **106**. In addition, clients **108**, **110** and **112** are also connected to network **102**. These clients, **108**, **110** and **112**, may be, for example, personal computers or network computers. For purposes of this application, a network computer is any computer coupled to a network that receives a program or other application from another computer coupled to the network. In the depicted example, server **104** is a logically partitioned platform and provides data, such as boot files, operating system images and applications, to clients **108-112**. Hardware system console **150** may be a laptop computer and is used to display messages to an

operator from each operating system image running on server **104**, as well as to send input information, received from the operator, to server **104**. Clients **108**, **110** and **112** are clients to server **104**. Distributed data processing system **100** may include additional servers, clients, and other devices not shown. Distributed data processing system **100** also includes printers **114**, **116** and **118**. A client, such as client **110**, may print directly to printer **114**. Clients such as client **108** and client **112** do not have directly attached printers. These clients may print to printer **116**, which is attached to server **104**, or to printer **118**, which is a network printer that does not require connection to a computer for printing documents. Client **110**, alternatively, may print to printer **116** or printer **118**, depending on the printer type and the document requirements.

In the depicted example, distributed data processing system **100** is the Internet, with network **102** representing a worldwide collection of networks and gateways that use the TCP/IP suite of protocols to communicate with one another. At the heart of the Internet is a backbone of high-speed data communication lines between major nodes or host computers consisting of thousands of commercial, government, education, and other computer systems that route data and messages. Of course, distributed data processing system **100** also may be implemented as a number of different types of networks such as, for example, an intranet or a local area network.

Figure 1 is intended as an example and not as an architectural limitation for the processes of the present

invention.

With reference now to **Figure 2**, a block diagram of a data processing system in accordance with the present invention is illustrated. Data processing system **200** is an example of a hardware system console, such as hardware system console **150** depicted in **Figure 1**. Data processing system **200** employs a peripheral component interconnect (PCI) local bus architecture. Although the depicted example employs a PCI bus, other bus architectures, such as Micro Channel and ISA, may be used. Processor **202** and main memory **204** are connected to PCI local bus **206** through PCI bridge **208**. PCI bridge **208** may also include an integrated memory controller and cache memory for processor **202**. Additional connections to PCI local bus **206** may be made through direct component interconnection or through add-in boards. In the depicted example, local area network (LAN) adapter **210**, SCSI host bus adapter **212**, and expansion bus interface **214** are connected to PCI local bus **206** by a direct component connection. In contrast, audio adapter **216**, graphics adapter **218**, and audio/video adapter (A/V) **219** are connected to PCI local bus **206** by add-in boards inserted into expansion slots. Expansion bus interface **214** provides a connection for a keyboard and mouse adapter **220**, modem **222**, and additional memory **224**. In the depicted example, SCSI host bus adapter **212** provides a connection for hard disk drive **226**, tape drive **228**, CD-ROM drive **230**, and digital video disc read only memory drive (DVD-ROM) **232.** Typical PCI local bus implementations will support three or four PCI expansion slots or add-in connectors.

An operating system runs on processor **202** and is used to coordinate and provide control of various components within data processing system **200** in **Figure 2**. The operating system may be a commercially available operating system, such as OS/2, which is available from International Business Machines Corporation. "OS/2" is a trademark of International Business Machines Corporation. An object-oriented programming system, such as Java, may run in conjunction with the operating system, providing calls to the operating system from Java programs or applications executing on data processing system **200**. Instructions for the operating system, the object-oriented operating system, and applications or programs are located on a storage device, such as hard disk drive **226**, and may be loaded into main memory **204** for execution by processor **202**.

Those of ordinary skill in the art will appreciate that the hardware in **Figure 2** may vary depending on the implementation. For example, other peripheral devices, such as optical disk drives and the like, may be used in addition to or in place of the hardware depicted in **Figure 2**. The depicted example is not meant to imply architectural limitations with respect to the present invention. For example, the processes of the present invention may be applied to multiprocessor data processing systems.

With reference now to **Figure 3**, a block diagram of a data processing system, which may be implemented as a logically partitioned server, such as server **104** in **Figure 1**, is depicted in accordance with the present invention. Data processing system **300** may be a symmetric

multiprocessor (SMP) system including a plurality of processors **301**, **302**, **303**, and **304** connected to system bus **306**. For example, data processing system **300** may be an IBM RS/6000, a product of International Business Machines Corporation in Armonk, New York. Alternatively, a single processor system may be employed. Also connected to system bus **306** is memory controller/cache **308**, which provides an interface to a plurality of local memories **360-363**. I/O bus bridge **310** is connected to system bus **306** and provides an interface to I/O bus **312**. Memory controller/cache **308** and I/O bus bridge **310** may be integrated as depicted.

Data processing system **300** is a logically partitioned data processing system. Thus, data processing system **300** may have multiple heterogeneous operating systems (or multiple instances of a single operating system) running simultaneously. Each of these multiple operating systems may have any number of software programs executing within in it. Data processing system **300** is logically partitioned such that different I/O adapters **320-321**, **328-329**, **336-337**, and **346-347** may be assigned to different logical partitions.

Thus, for example, suppose data processing system **300** is divided into three logical partitions, P1, P2, and P3. Each of I/O adapters **320-321**, **328-329**, and **336-337**, each of processors **301-304**, and each of local memories **360-364** is assigned to one of the three partitions. For example, processor **301**, memory **360**, and I/O adapters **320**, **328**, and **329** may be assigned to logical partition P1; processors **302-303**, memory **361**, and I/O adapters **321** and

337 may be assigned to partition P2; and processor **304**, memories **362-363**, and I/O adapters **336** and **346-347** may be assigned to logical partition P3.

Each operating system executing within data processing system **300** is assigned to a different logical partition. Thus, each operating system executing within data processing system **300** may access only those I/O units that are within its logical partition. Thus, for example, one instance of the Advanced Interactive Executive (AIX) operating system may be executing within partition P1, a second instance (image) of the AIX operating system may be executing within partition P2, and a Windows 2000™ operating system may be operating within logical partition P1. Windows 2000 is a product and trademark of Microsoft Corporation of Redmond, Washington.

Peripheral component interconnect (PCI) Host bridge **314** connected to I/O bus **312** provides an interface to PCI local bus **315**. A number of Terminal Bridges **316-317** may be connected to PCI bus **315**. Typical PCI bus implementations will support four Terminal Bridges for providing expansion slots or add-in connectors. Each of Terminal Bridges **316-317** is connected to a PCI/I/O Adapter **320-321** through a PCI Bus **318-319**. Each I/O Adapter **320-321** provides an interface between data processing system **300** and input/output devices such as, for example, other network computers, which are clients to server **300**. Only a single I/O adapter **320-321** may be connected to each terminal bridge **316-317**. Each of terminal bridges **316-317** is configured to prevent the

propagation of errors up into the PCI Host Bridge **314** and into higher levels of data processing system **300**. By doing so, an error received by any of terminal bridges **316-317** is isolated from the shared buses **315** and **312** of the other I/O adapters **321**, **328-329**, and **336-337** that may be in different partitions. Therefore, an error occurring within an I/O device in one partition is not "seen" by the operating system of another partition. Thus, the integrity of the operating system in one partition is not effected by an error occurring in another logical partition. Without such isolation of errors, an error occurring within an I/O device of one partition may cause the operating systems or application programs of another partition to cease to operate or to cease to operate correctly.

Additional PCI host bridges **322**, **330**, and **340** provide interfaces for additional PCI buses **323**, **331**, and **341**. Each of additional PCI buses **323**, **331**, and **341** are connected to a plurality of terminal bridges **324-325**, **332-333**, and **342-343** which are each connected to a PCI I/O adapter **328-329**, **336-337**, and **346-347** by a PCI bus **326-327**, **334-335**, and **344-345**. Thus, additional I/O devices, such as, for example, modems or network adapters may be supported through each of PCI I/O adapters **328-329, 336-337**, and **346-347**. In this manner, server **300** allows connections to multiple network computers. A memory mapped graphics adapter **348** and hard disk **350** may also be connected to I/O bus **312** as depicted, either directly or indirectly. Hard disk **350** may be logically partitioned between various partitions without the need

for additional hard disks. However, additional hard disks may be utilized if desired.

Those of ordinary skill in the art will appreciate that the hardware depicted in **Figure 3** may vary. For example, other peripheral devices, such as optical disk drives and the like, also may be used in addition to or in place of the hardware depicted. The depicted example is not meant to imply architectural limitations with respect to the present invention.

With reference now to **Figure 4**, a block diagram illustrating a prior art logically partitioned platform is depicted in accordance with the present invention. Logically partitioned platform **400** is an example of a platform that, in prior art systems, may have been implemented as server **104** in **Figure 1**. Logically partitioned platform **400** includes partitioned hardware **430**, shared single hardware **420**, and operating systems **402-408**. Operating systems **402-408** may be multiple copies of a single operating system or multiple heterogeneous operating systems simultaneously run on platform **400**.

Partitioned hardware **430** includes a plurality of processors **432-438**, a plurality of system memory units **440-446**, a plurality of input/output (I/O) adapters **448-462**, and a storage unit **470**. Each of the processors **442-448**, memory units **440-446**, and I/O adapters **448-462** may be assigned to one of multiple partitions within logically partitioned platform **400**, each of which corresponds to one of operating systems **402-408**.

Shared single hardware unit **420** includes console **422**

and operator panel **424**. Shared single hardware unit **420** may also include other shared devices not depicted in **Figure 4**. Console **422** typically includes a display and data entry device such as a keyboard. Console **422** allows an operator to respond to and correct errors displayed on operator panel **424**. Operator panel **424** is typically a panel display, such as an LCD display, on the front of the physical chassis of the server in which text messages are displayed alerting an operator of potential problems within platform **400** or within a particular OS **402-408** running on platform **400**.

Each operating system image **402-408** must share access to resources within shared single hardware **420**. Therefore, some of the benefits of a logically partitioned platform are lost, since each partition may access and change the contents of shared resources, thus affecting other partitions within the platform. One alternative to allowing each partition to share access to shared single hardware **420** is to duplicate these hardware devices and have a separate operator panel and console for each partition. However, such a solution is bulky and, often cost prohibitive.

With reference now to **Figure 5**, a block diagram of an exemplary logically partitioned platform is depicted in which the present invention may be implemented. The hardware in logically partitioned platform **500** may be implemented as, for example, server **200** in **Figure 2**. Logically partitioned platform **500** is similar to logically partitioned platform **400** in **Figure 4**. However, Hypervisor **510**, implemented as firmware, has been added.

Firmware is "hard software" stored in a memory chip that holds its content without electrical power, such as, for example, read-only memory (ROM), programmable ROM (PROM), erasable programmable ROM (EPROM), electrically erasable programmable ROM (EEPROM), and non-volatile random access memory (non-volatile RAM).

Hypervisor **510** provides the OS images **402-408** running in multiple logical partitions each a virtual copy of a console and operator panel. The interface to the console is changed from an asynchronous teletype port device driver, as in the prior art, to a set of hypervisor firmware calls that emulate a port device driver. The hypervisor **510** encapsulates the data from the various OS images onto a message stream that is transferred to a computer **580**, known as a hardware system console **580**.

Hardware system console **580** is connected directly to logically partitioned platform **500** as illustrated in **Figure 5**, or may be connected to logically partitioned platform through a network, such as, for example, network **102** in **Figure 1**. Hardware system console **580** may be, for example, a desktop or laptop computer, and may be implemented as data processing system **200** in **Figure 2**. Hardware system console **580** decodes the message stream and displays the information from the various OS images **402-408** in separate windows, at least one per OS image. Similarly, keyboard input information from the operator is packaged by the hardware system console, sent to logically partitioned platform **500** where it is decoded and delivered to the appropriate OS image via the hypervisor **510** emulated port device driver associated

with the then active window on the hardware system console **580**.

Those of ordinary skill in the art will appreciate that the hardware and software depicted in **Figure 5** may vary. For example, more or fewer processors and/or more or fewer operating system images may be used than those depicted in **Figure 5**. The depicted example is not meant to imply architectural limitations with respect to the present invention.

With reference now to **Figures 6A-6B**, high-level flowcharts illustrating exemplary processes, performed for example, in hypervisor **510**, for emulating a console and operator platform is depicted in accordance with the present invention. The operating systems, such as, for example, OS **402-408** in **Figure 4**, call the hypervisor through a single entry point. One thread of execution illustrated in **Figure 6A** gets data from a per partition buffer and sends data to the hardware system console while a separate thread of execution illustrated in **Figure 6B** receives data from the hardware system console.

In the first thread of execution depicted in **Figure 6A**, the hypervisor receives a request from an operating system to get or send data (step **601**). The hypervisor determines whether the request is a request to send or get data (step **602**). If the request is a request to send data, then the hypervisor determines from which OS image (partition) the received data originated (step **604**). The received data is then encapsulated onto a message stream (step **606**). The encapsulated data includes the message or information received from the OS as well as the

identity of the OS. The hypervisor then sends the message stream to the hardware system console (step **608**).

If the request is a request to get data, then the hypervisor determines which OS partition requested the data (step **610**). Each partition is assigned a data buffer for storing data received from the hardware system console until retrieved by the partition's OS. Thus, the hypervisor determines whether the requesting partition's data buffer is empty (step **612**). If the data buffers for the requesting partition is empty, then the hypervisor sends a NULL message to the requesting OS image indicating that there is no data from the hardware system console for the OS image to receive (step **616**). If the data buffer is not empty, then the message data from the partition's data buffer is sent to the requesting OS image (step **614**).

In the second thread of execution depicted in **Figure 6B**, the hypervisor receives and decodes data from the hardware system console (step **618**). The hypervisor then places the decoded data into the buffer corresponding to the appropriate partition such that it may be retrieved and sent to the appropriate partition's OS image upon request (step **620**).

With reference now to **Figure 7**, a high level flowchart illustrating an exemplary process on a hardware system console for presenting the information from the various OS images to an operator is depicted in accordance with the present invention. To begin, the hardware system console receives a message stream from the hypervisor (step **702**). The hardware system console decodes the message stream (step **704**) and determines to

which OS image the received data corresponds (step **706**). Next, the hardware system console determines which window within the display corresponds to the determined OS image (step **708**) and displays the received data to an operator in the window corresponding to the proper OS image (step **710**).

With reference now to **Figure 8**, a high level flowchart illustrating an exemplary process on a hardware system console for sending messages to various ones of multiple OS images running on a logically partitioned platform is depicted in accordance with the present invention. To begin, the hardware system console receives input from an operator from an input device, such as, for example, a keyboard (step **802**). The hardware system console then determines which OS image corresponds to the active window from which the input was received (step **804**). The data input, along with the OS image it corresponds with, are encapsulated into a message stream (step **806**). The message stream is then sent to the hypervisor (step **808**).

It is important to note that while the present invention has been described in the context of a fully functioning data processing system, those of ordinary skill in the art will appreciate that the processes of the present invention are capable of being distributed in the form of a computer readable medium of instructions and a variety of forms and that the present invention applies equally regardless of the particular type of signal bearing media actually used to carry out the distribution. Examples of computer readable media include recordable-type media such as a floppy disc, a

hard disk drive, a RAM, and CD-ROMs and transmission-type media such as digital and analog communications links.

The description of the present invention has been presented for purposes of illustration and description, but is not intended to be exhaustive or limited to the invention in the form disclosed. Many modifications and variations will be apparent to those of ordinary skill in the art. The embodiment was chosen and described in order to best explain the principles of the invention, the practical application, and to enable others of ordinary skill in the art to understand the invention for various embodiments with various modifications as are suited to the particular use contemplated.

CLAIMS:

What is claimed is:

1. A logically partitioned data processing system, comprising:

a plurality of logical partitions;

a plurality of operating systems, each assigned to a separate one of the plurality of logical partitions;

a plurality of assignable resources, wherein each of the plurality of assignable resources is assigned to one of the plurality of logical partitions;

a hypervisor, wherein the hypervisor emulates shared resources and provides a virtual copy of the shared resources to each of the plurality of logical partitions.

2. The logically partitioned data processing system as recited in claim 1, wherein the shared resources comprise an operator panel.

3. The logically partitioned data processing system as recited in claim 1, wherein the shared resources comprise a system console.

4. The logically partitioned data processing system as recited in claim 1, wherein the hypervisor receives a system message from one of the plurality of operating system images, appends an operating system identity to the message to produce a new message, and sends the new

message to an external data processing system.

5. The logically partitioned data processing system as recited in claim 1, wherein instructions for executing the hypervisor are contained within firmware.

6. The logically partitioned data processing system as recited in claim 5, wherein the firmware comprises a read-only memory.

7. The logically partitioned data processing system as recited in claim 5, wherein the firmware comprises a programmable read-only memory.

8. The logically partitioned data processing system as recited in claim 5, wherein the firmware comprises an erasable programmable read-only memory.

9. The logically partitioned data processing system as recited in claim 5, wherein the firmware comprises an electrically erasable programmable read-only memory.

10. The logically partitioned data processing system as recited in claim 5, wherein the firmware comprises a non-volatile random access memory.

11. A method of providing separate copies of shared resources to each of multiple partitions within a data processing system, the method comprising:

receiving, at a hypervisor, a message from a one of a plurality of operating system images, executing within

the data processing system, intended for a shared resource;

determining an identity of the one of the plurality of operating system images;

encoding the message and the identity into a new message; and

transmitting the new message to an external data processing system for presentation to a user.

12. The method as recited in claim 11, wherein the shared resource is an operator panel.

13. The method as recited in claim 11, wherein the shared resource is a system console.

14. The method as recited in claim 11, further comprising:

receiving external data from the external data processing system;

decoding the external data to determine an input, an identity of the shared resource, and an intended one of the plurality of operating system images; and

transmitting the input to the intended one of the plurality of operating system images with an indication the identity of the shared resource from which the input corresponds.

15. A computer program product for providing separate copies of shared resources to each of multiple partitions within a data processing system, the computer program product comprising:

first instructions for receiving, at a hypervisor, a message from a one of a plurality of operating system images, executing within the data processing system, intended for a shared resource;

second instructions for determining an identity of the one of the plurality of operating system images;

third instructions for encoding the message and the identity into a new message; and

fourth instructions for transmitting the new message to an external data processing system for presentation to a user.

16. The computer program product as recited in claim 15, wherein the shared resource is an operator panel.

17. The computer program product as recited in claim 15, wherein the shared resource is a system console.

18. The computer program product as recited in claim 15, further comprising:

fifth instructions for receiving external data from the external data processing system;

sixth instructions for decoding the external data to determine an input, an identity of the shared resource, and an intended one of the plurality of operating system images; and

seventh instructions for transmitting the input to the intended one of the plurality of operating system images with an indication the identity of the shared resource from which the input corresponds.

19. A system for providing separate copies of shared resources to each of multiple partitions within a data processing system, the system comprising:

first means for receiving, at a hypervisor, an message from a one of a plurality of operating system images, executing within the data processing system, intended for a shared resource;

second means for determining an identity of the one of the plurality of operating system images;

third means for encoding the message and the identity into a new message; and

fourth means for transmitting the new message to an external data processing system for presentation to a user.

20. The system as recited in claim 19, wherein the shared resource is an operator panel.

21. The system as recited in claim 19, wherein the shared resource is a system console.

22. The system as recited in claim 19, further comprising:

fifth means for receiving external data from the external data processing system;

sixth means for decoding the external data to determine an input, an identity of the shared resource, and an intended one of the plurality of operating system images; and

seventh means for transmitting the input to the intended one of the plurality of operating system images with an

indication the identity of the shared resource from which the input corresponds.

23. A system for partitioning shared resources, the system comprising:

a first data processing system comprising:

a plurality of partitions each corresponding to separate one of a plurality of operating system images;

a plurality of assignable resources; and

a hypervisor for providing each partition a separate one of a shared system resource;

a second data processing system coupled to the first data processing system, wherein the second data processing system receives a message from the hypervisor, wherein the message indicates to which of the plurality of operating system images the message belong, and wherein the second data processing system displays the message to a user with an indication of the operating system image corresponding to the message.

24. The system as recited in claim 23, wherein the data processing system, responsive to operator input for a specified one of the plurality of operating system images, sends encapsulated data, comprising the operator input and an indication of the corresponding operating system image, to the hypervisor, and wherein the hypervisor decodes the encapsulated data and sends the operator input to the corresponding operating system image.

ABSTRACT OF THE DISCLOSURE

HYPERVISOR VIRTUALIZATION OF OS CONSOLE AND OPERATOR PANEL

A logically partitioned data processing system in which shared resources are emulated to provide each partition a separate copy of the shared resource is provided. In one embodiment, the logically partitioned data processing system includes a plurality of logical partitions, a plurality of operating systems executing within the data processing system and a plurality of assignable resources. Each of the plurality of operating systems is assigned to a separate one of the plurality of logical partitions, such that no more than one operating system is assigned to any given logical partition. Each of the plurality of assignable resources is assigned to a single one of the plurality of logical partitions. The logically partitioned data processing system also includes a hypervisor. The hypervisor emulates shared resources, such as an operator panel and a system console, and provides a virtual copy of these shared resources to each of the plurality of logical partitions.

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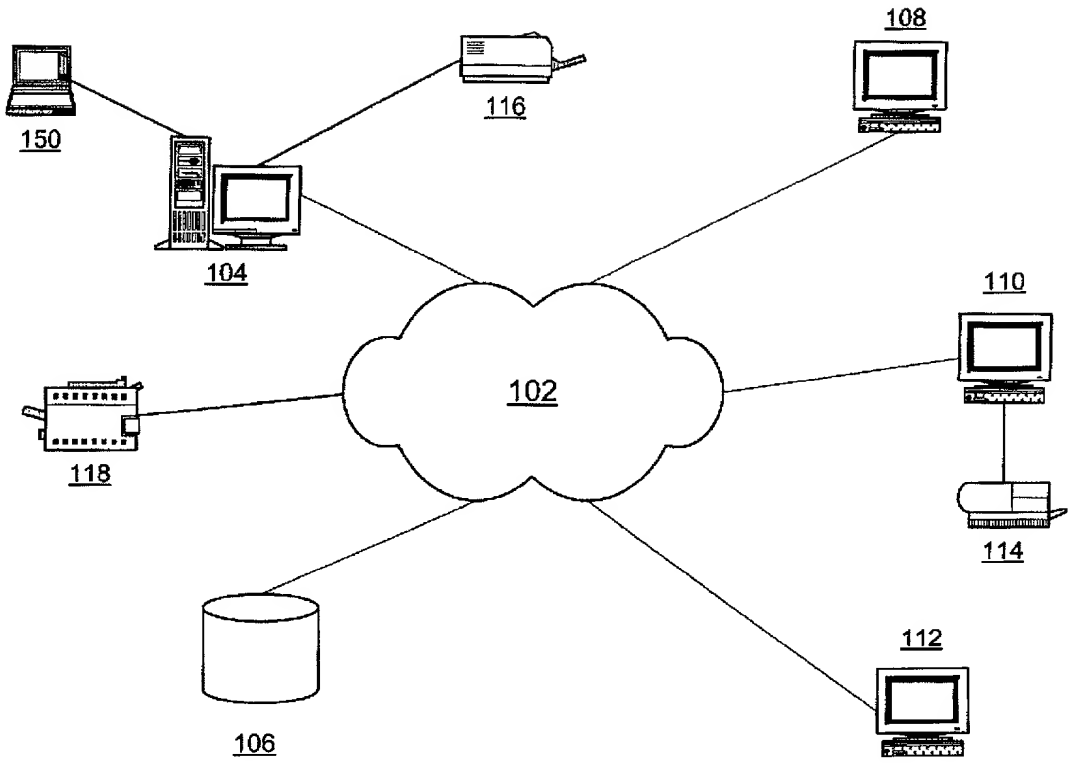
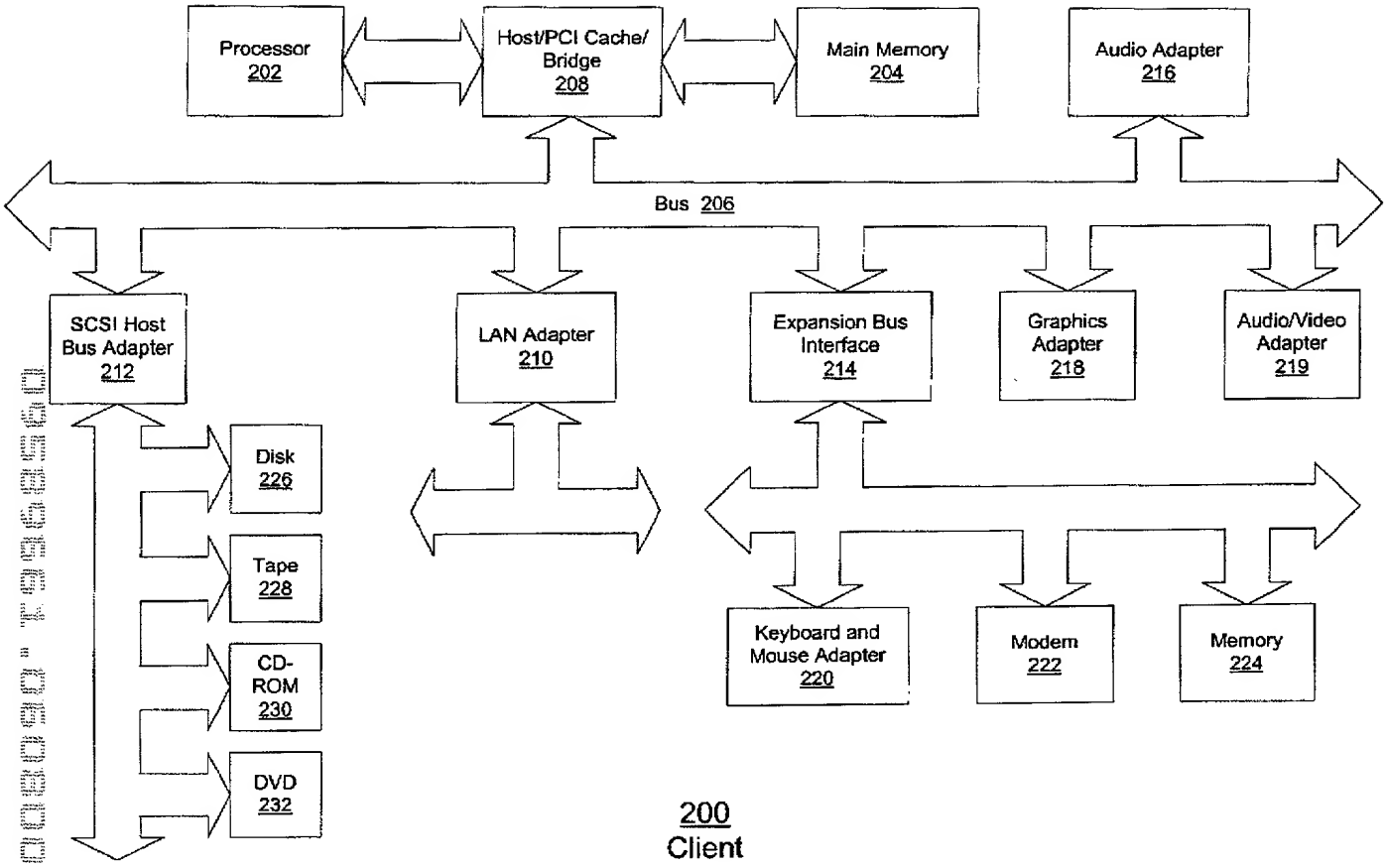


100

Network

Figure 1



200
Client

Figure 2

AUS990938US1

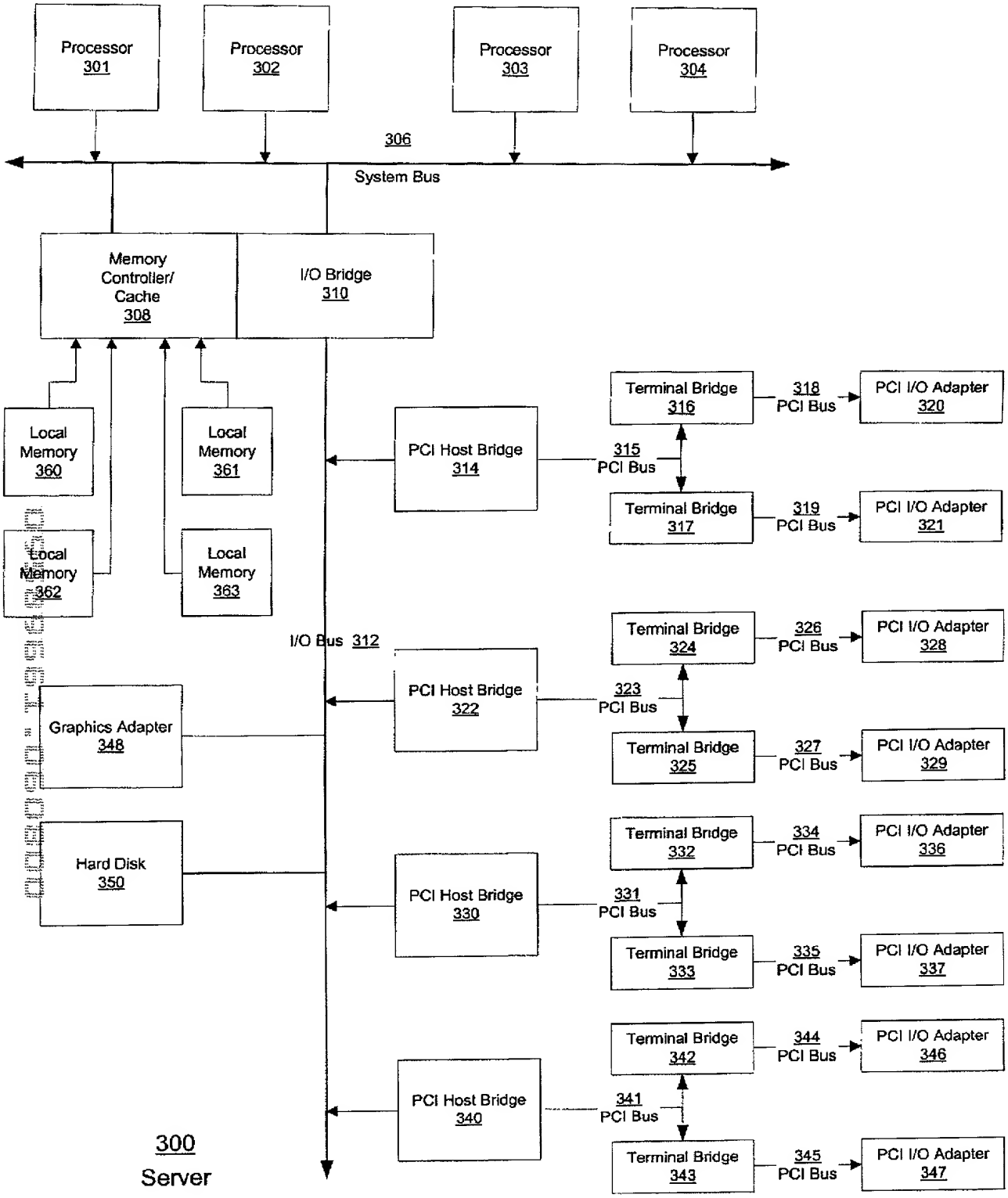
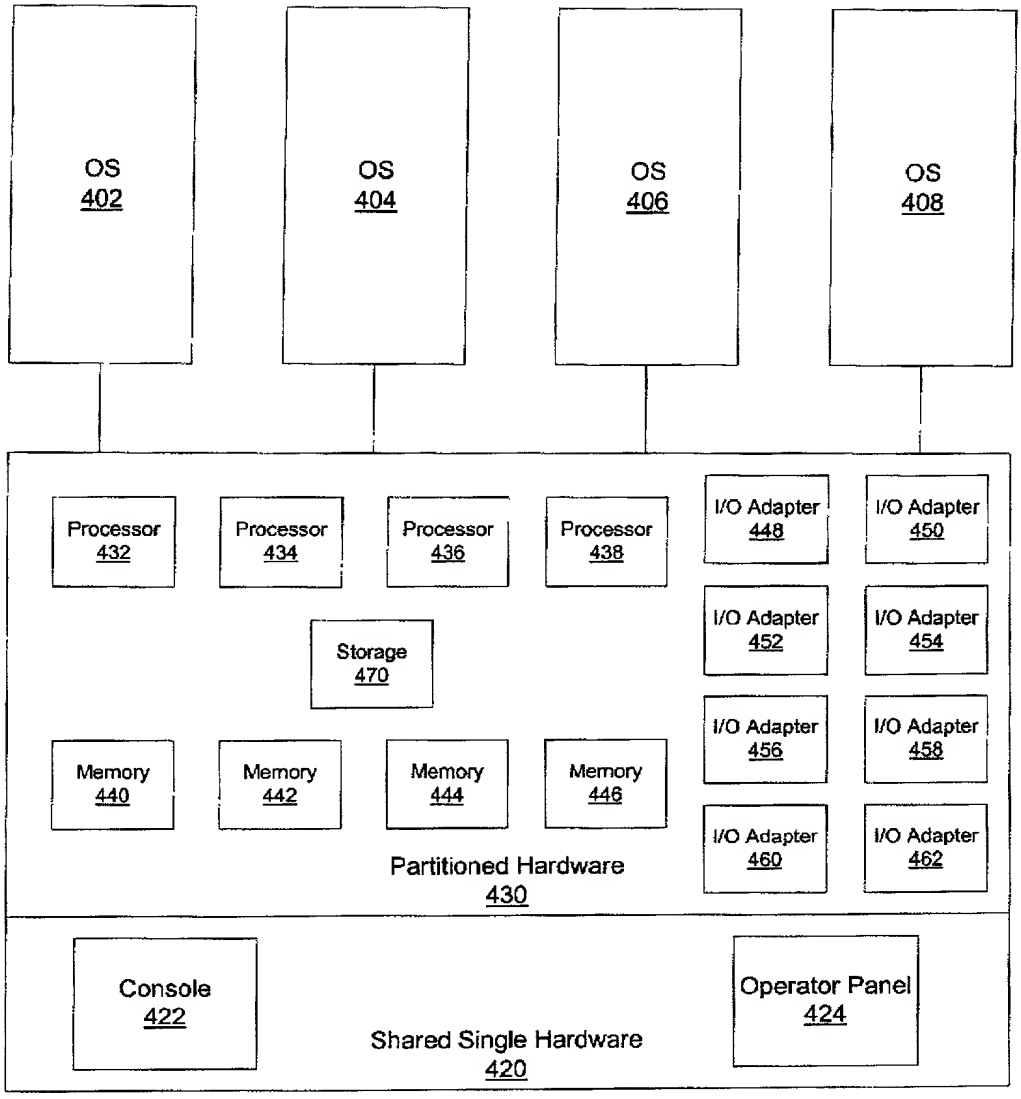


Figure 3

AUS990938US1



400
Logically Partitioned Platform

(Prior Art)

Figure 4

AUS990938US1

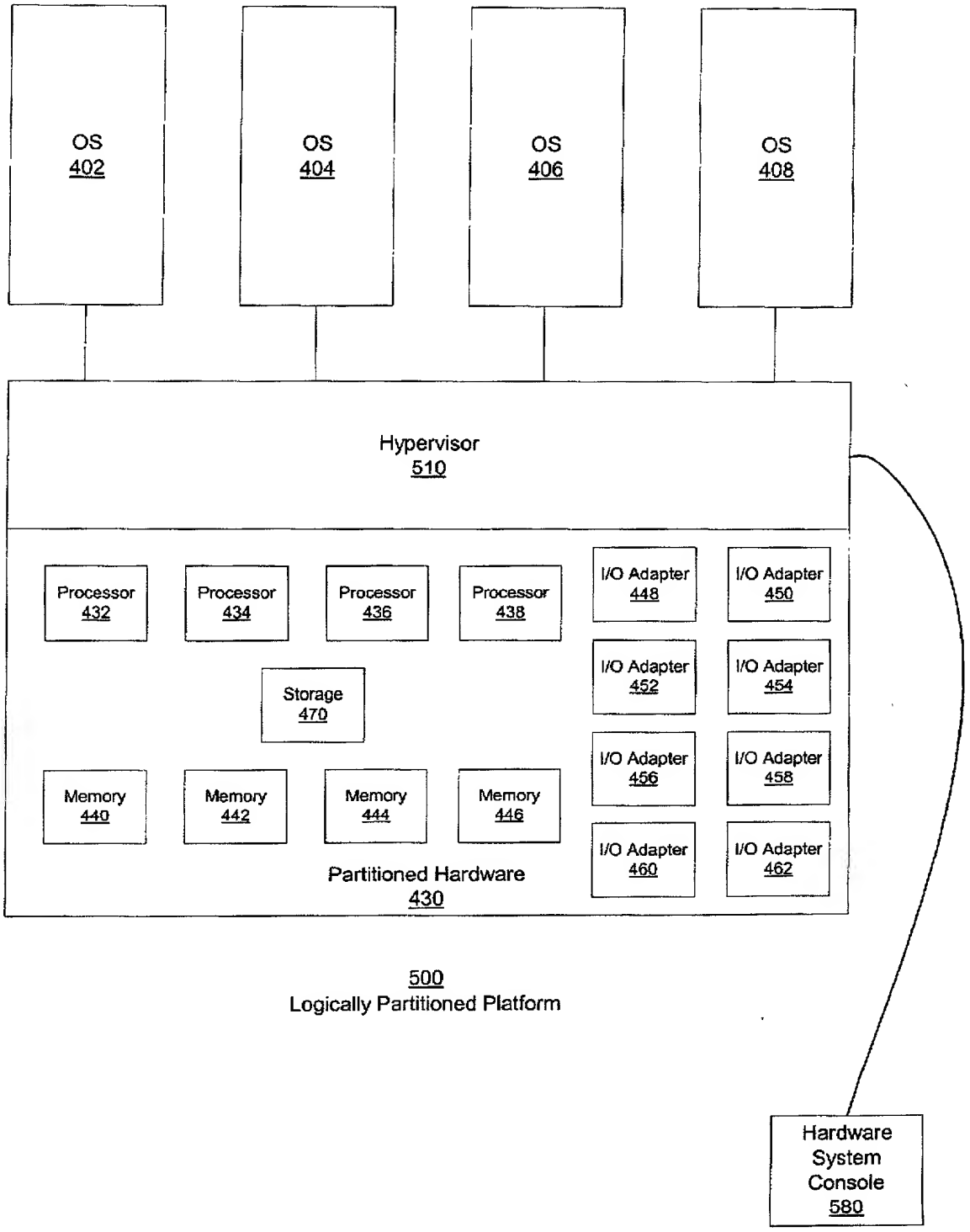


Figure 5

AUS990938US1

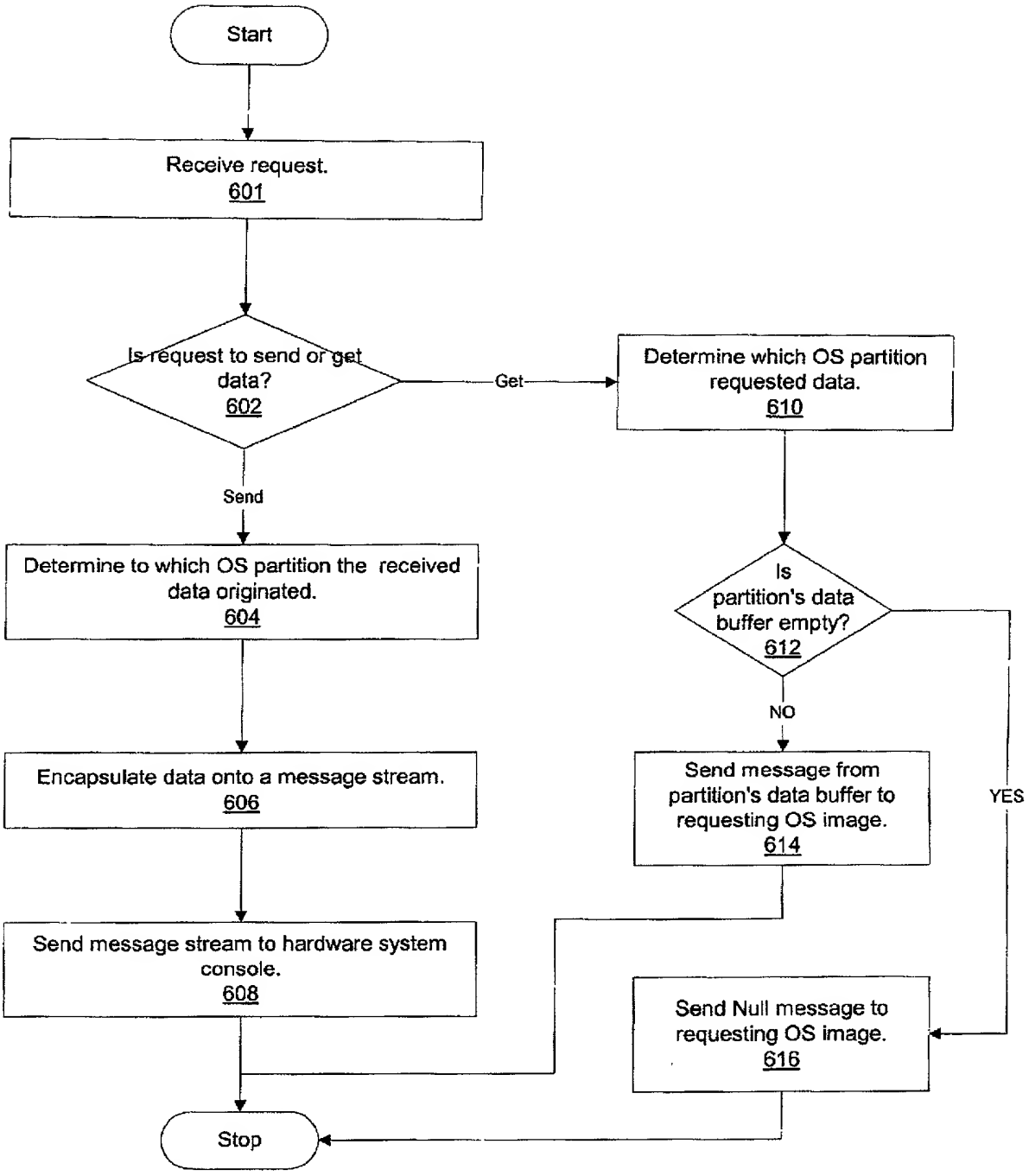


Figure 6A

AUS990938US1

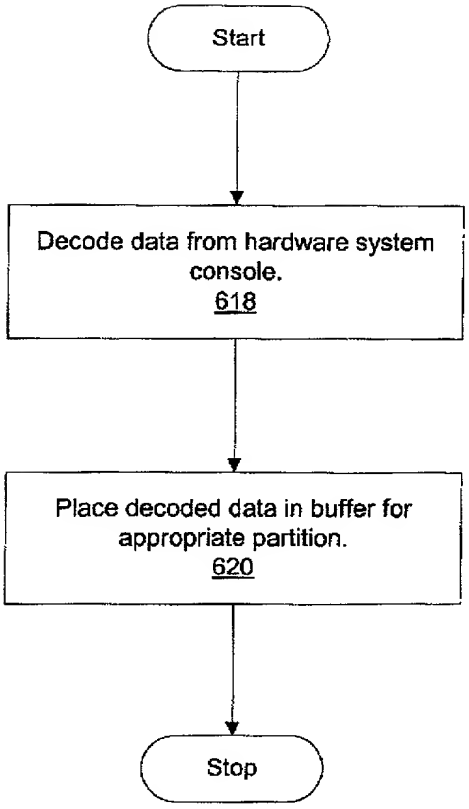

Start
Decode data from hardware system console.
618
Place decoded data in buffer for appropriate partition.
620
Stop

Figure 6B

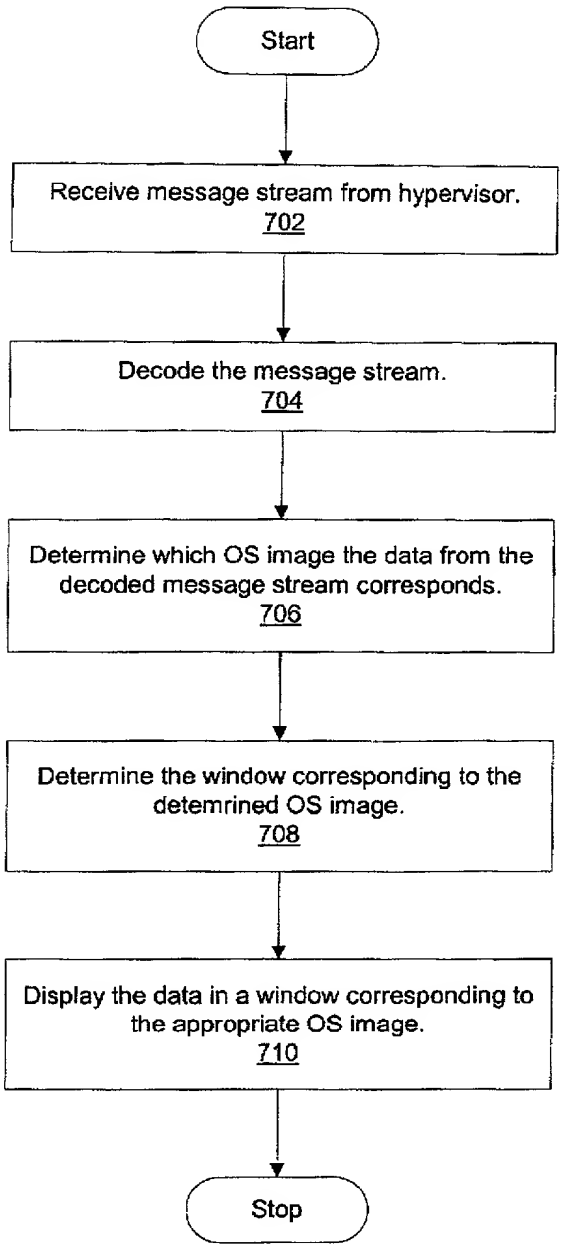


Figure 7

AUS990938US1

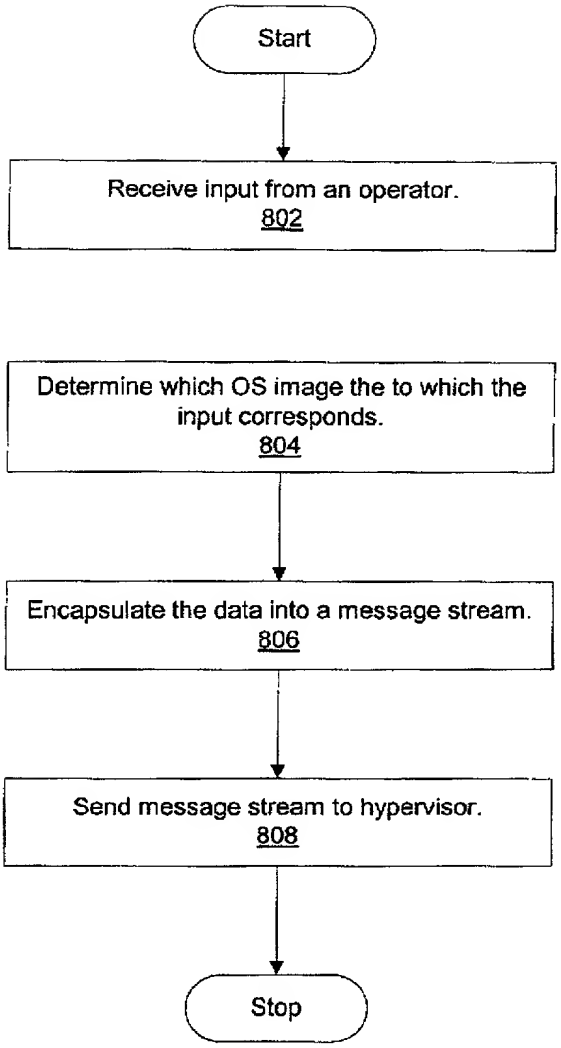


Figure 8

AUS990938US1

DOCKET NUMBER: AUS990938US1

DECLARATION AND POWER OF ATTORNEY FOR

PATENT APPLICATION

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name;

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

HYPERVISOR VIRTUALIZATION OF OS CONSOLE AND OPERATOR PANEL
the specification of which (check one)

X is attached hereto.

__ was filed on _____
as Application Serial No._____
and was amended on _____
(if applicable)

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to the patentability of this application in accordance with Title 37, Code of Federal Regulations, §1.56.

I hereby claim foreign priority benefits under Title 35, United States Code, §119 of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which priority is claimed:

Prior Foreign Application(s):			Priority Claimed
_____ (Number)	_____ (Country)	_____ (Day/Month/Year)	___ Yes___ No

I hereby claim the benefit under Title 35, United States Code, §120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, §112, I acknowledge the duty to disclose information material to the patentability of this application as defined in Title 37, Code of Federal Regulations, §1.56 which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

_____ (Application Serial #)	_____ (Filing Date)	_____ (Status)

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

POWER OF ATTORNEY: As a named inventor, I hereby appoint the following attorneys and/or agents to prosecute this application and transact all business in the Patent and Trademark Office connected therewith.

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